# Algorithms on Graphs

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Day 2, session 1: Walks in graphs



STCS Vigyan Vidushi 2024



Paths vs walks:

Path: each vertex appears at most once

Walk: a vertex can appear multiple times



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Cycles vs circuits:

Cycle: each vertex appears at most once (except for first and last)

Circuit: a vertex can appear multiple times



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<u>Eulerian walk</u>: sequence of adjacent edges, such that each edge appears exactly once.



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Start from b, end at d

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Start from b, end at d



Start from d, end at c

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<u>Eulerian circuit</u>: sequence of adjacent edges, such that each edge appears exactly once, and begins and ends at the same vtx.





#### no Eulerian circuit

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<u>Theorem:</u> Graph G = (V,E) has an Eulerian ckt iff every vertex has even degree



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Eulerian ckt  $\Rightarrow$  even degree: easy

Even degree  $\Rightarrow$  Eulerian ckt: via algorithms



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Hierholzer's Algorithm:

from any vertex (say e),

keep following unused edges, until back to e

store vertices on this trail that have other edges



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repeat procedure for these vertices



е

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<u>Theorem:</u> Graph G = (V,E) has an Eulerian ckt iff every vertex has even degree



16 cards were in a fixed sequence:

1 **4** , 4 **4** , 8 **4** , 2 **4** , 3 **9** , 5 **4** , 3 **4** , 9 **4** , 2 **4** , 1 **4** , 7 **9** , 7 **4** , 1 **4** , 5 **4** , 4 **9** , 2 **9** 

16 cards were in a fixed sequence:

1 🛧 , 4 🛧 , 8 🛧 , 2 🛧 , 3 🤍 , 5 🛧 , 3 🛧 , 9 🔶 , 2 🔶 , 1 🛧 , 7 🧡 , 7 🛧 , 1 🔶 , 5 🔶 , 4 🧡 , 2 🧡

One (or more) straight cut only permutes the cards cyclically:



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1 🛧 , 4 🛧 , 8 🛧 , 2 🛧 , 3 🧡 , 5 🛧 , 3 🛧 , 9 🔶 , 2 🔶 , 1 🛧 , 7 🧡 , 7 🛧 , 1 🔶 , 5 🔶 , 4 🧡 , 2 🧡

One (or more) straight cut only permutes the cards cyclically:

A sequence of four cards is picked from the top:

But how would I know which card is on top?

Let's look at the sequence again:

#### 1 🛧 , 4 🛧 , 8 🛧 , 2 🛧 , 3 🧡 , 5 🛧 , 3 🛧 , 9 🔶 , 2 🔶 , 1 🛧 , 7 🧡 , 7 🛧 , 1 🔶 , 5 🔶 , 4 🧡 , 2 🧡

#### Let's look at the sequence again:



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Each 4-sequence of Bs and Rs appears exactly once in this sequence (considered cyclically)!

 $(2^4 = 16 \text{ such sequences}, \& 16 \text{ cards}, \text{hence exactly } 16 \text{ starting points})$ 

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When I asked all Red card holders to raise their hand, I got a sequence



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gives

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 $k = 2: \quad 0 \quad 0 \quad 1 \quad 1$   $k = 3: \quad 0 \quad 0 \quad 0 \quad 1 \quad 0 \quad 1 \quad 1 \quad 1$  $k = 4: \quad \mathbf{B} \quad \mathbf{B} \quad \mathbf{B} \quad \mathbf{B} \quad \mathbf{R} \quad \mathbf{B} \quad \mathbf{R} \quad$ 

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k = 2: 0 0 1 1 k = 3: 0 0 0 1 0 1 1 1 k = 4: B B B B R B B R R R R R R R R R0 0 0 0 1 0 1 0 1 1 0 1 0 1 1 1

But do such sequences always exist? How do we find them?

#### k = 3: 0 0 0 1 0 1 1 1









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$$k = 3$$
: 0 0 0 1 0 1 1 1  
Each vtx has  
in-degree = out-degree = 2

8 edges







000 k = 3: 0 0 0 1 0 1 1 1 To generate a k-window de Bruijn 00 100 001 sequence: 010 - draw graph with  $2^{k-1}$  nodes, 10 01 each node with length k-1101 binary string - add edges (2 incoming and 011 110 11 outgoing for each node), each edge gives length k string 111 - find an Eulerian walk

### **Other Uses**

de Bruijn sequences useful for many applications:

- robotics, for robots to decide where they are
- encoding and decoding
- genome reconstruction

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Thanks for listening!